

# *LINKED STRUCTURES IN JAVA*

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# Say You Want To Design a Container

- That can should be able to hold a variable amount of data
  - Where “variable” means not known at compile time
- Three ways
- First, you could use an existing container
- Or, you could use an array (since can be sized at runtime)
- Or, you can build a linked structure
  - This is what we’ll consider today!

# Linked Structures

- Are always declared like this:

```
public class LinkedList <T extends Comparable <T>> {  
  
    private T myData;  
    private LinkedList <T> next;  
  
    public LinkedList () {  
        next = null;  
        myData = null;  
    }  
}
```

- By “always” I mean that they always have this recursive structure
- Where you have a class that contains a reference to an object of its own type
- But naturally, are infinite variations on basic idea!

# How They Grow

- Can hold a variable amount of data ‘cause you can add new objects to the chain
- How to write an “insert (T insertMe)” method?

# How They Grow

```
public class LinkedList <T extends Comparable <T>> {  
  
    // convention: myData in root is null  
    // myData in everyone else is valid  
    // makes it possible to do "direct" removals  
    private T myData;  
    private LinkedList <T> next;  
  
    private LinkedList (T data, LinkedList <T> nextOne) {  
        myData = data;  
        next = nextOne;  
    }  
  
    public void insert (T insertMe) {  
        next = new LinkedList <T> (insertMe, next);  
    }  
}
```

# Why Does This Follow Recursion in the Syl?

- Because linked structures are recursive in nature
- Often easiest to write methods for them recursively
- Consider writing a “remove (T removeMe)” method...

# Why Does This Follow Recursion in the Syl?

```
public class LinkedList <T extends Comparable <T>> {  
  
    private T myData;  
    private LinkedList <T> next;  
  
    public T remove (T removeMe) {  
        if (next == null) {  
            return null;  
        } else if (next.myData.compareTo (removeMe) == 0) {  
            T returnVal = next.myData;  
            next = next.next;  
            return returnVal  
        } else {  
            return next.remove (removeMe);  
        }  
    }  
}
```

# How About Printing In Order of Insertion?

- Note that the last item added is at the front of the list
- So we want to print from the back to the front

# How About Printing In Order of Insertion?

```
public class LinkedList <T extends Comparable <T>> {  
  
    private T myData;  
    private LinkedList <T> next;  
  
    // since myData is null at root, we always print  
    // the list starting at "next"  
    public void print () {  
        if (next == null) {  
            return;  
        } else {  
            next.print ();  
            System.out.println (next.myData);  
        }  
    }  
}
```

# Easy, Right?

- Well, linked structures can be of arbitrary complexity
- A common (more complex) linked structure: BST
- Each BST node has two children: left and right
- When inserting, maintain the invariant:
  - Data in root is no smaller than everything in left subtree
  - Data in root is less than everything in right subtree
- Allows fast,  $\log(n)$  lookups if “balanced”
  - At every node, depth of left subtree and right subtree differs by at most a constant
- Are many flavors of BSTs

# Insertion Into a Simple BST

```
public class BST <T extends Comparable <T>> {  
  
    private T myData;  
    private BST <T> leftSubtree;  
    private BST <T> rightSubtree;  
  
    public BST () {  
        myData = null;  
        leftSubtree = rightSubtree = null;  
    }  
  
    public void insert (T insertMe) {  
        if (myData == null) {  
            myData = insertMe;  
        } else {...}  
    }  
}
```

# Insertion Into a Simple BST

```
public class BST <T extends Comparable <T>> {  
  
    public void insert (T insertMe) {  
        if (myData == null) {  
            myData = insertMe;  
        } else if (myData.compareTo (insertMe) >= 0) {  
            if (leftSubtree == null) {  
                leftSubtree = new BST <T> ();  
            }  
            leftSubtree.insert (insertMe);  
        } else {  
            if (rightSubtree == null) {  
                rightSubtree = new BST <T> ();  
            }  
            rightSubtree.insert (insertMe);  
        }  
    }  
}
```

# How To Search?

- Just start at the root, and recurse down the tree

# Searching a Simple BST

```
public boolean isThere (T findMe) {
    if (myData == null) {
        return false;
    } else if (myData.compareTo (findMe) == 0) {
        return true;
    } else if (myData.compareTo (findMe) > 0) {
        return (leftSubtree != null) &&
            leftSubtree.isThere (findMe);
    } else {
        return (rightSubtree != null) &&
            rightSubtree.isThere (findMe);
    }
}
```

- Note that this does not crash ‘cause of “short circuiting”
- That is, if the first part of an “and” evals to false, second part is ignored

# So, How To Keep a BST Balanced?

- Our simple BST will only be balanced w. random insert order
- What if the insert order is not random?
  - Well, can “hash” the inserted objects then compare on hashed vals instead
  - Pros and cons?
- Many classic BST variants are “self balancing”
  - AVL trees
  - Red/black trees
  - All have intricate, challenging algorithms! Take a look!
- Won't cover in this class, since not that useful for our doc system
  - Instead, we'll cover another type of linked tree structure in depth
  - Called a “B-Tree”
  - Starting next time!

Questions?